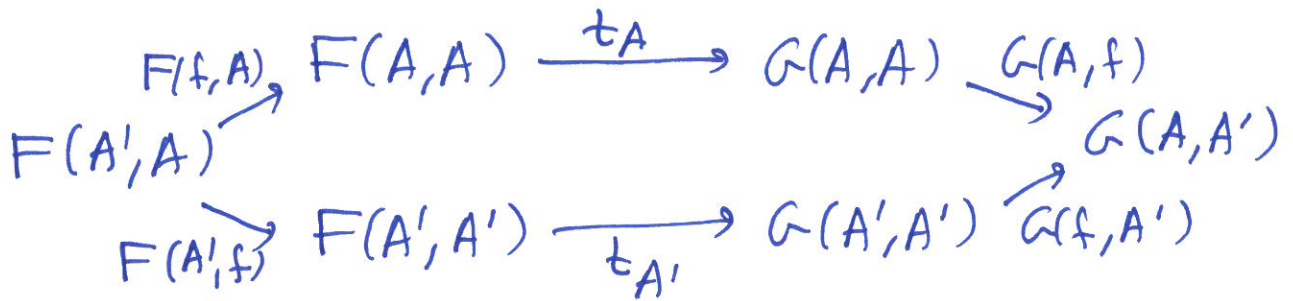


ENRICHED
DINATURAL
NUMBERS

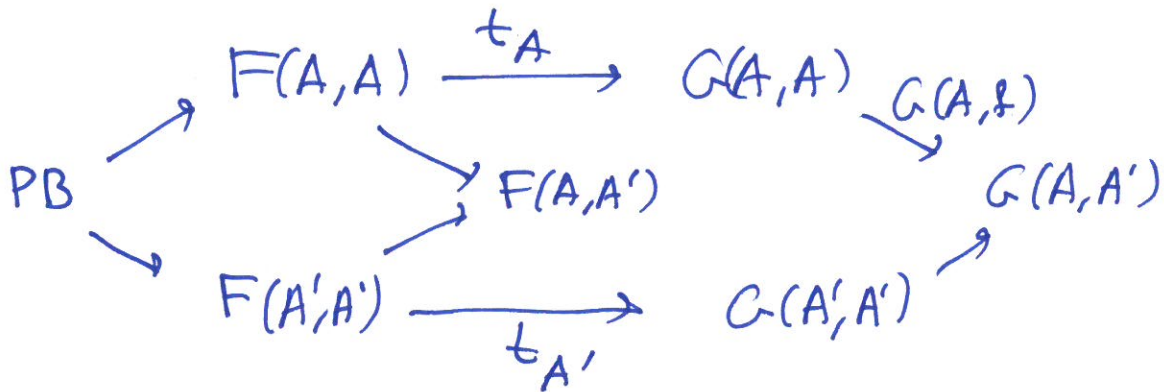
FEB. 7, 1995

For $F, G: \underline{A}^{op} \times \underline{A} \rightarrow \underline{B}$

DINATURAL TRANSFS:



BARR DINATURALS:



THEOREM: STRONG BARR DINATS $\text{Hom}_{\underline{A}} \Rightarrow \text{Hom}_{\underline{A}}^L$
 $\cong \frac{\text{STRONG BARR DINATS}}{L \rightarrow \underline{N}}$

$\text{Hom}_{\underline{A}}^L : \underline{A}^{op} \times \underline{A} \rightarrow \underline{\text{Set}} ; (A, A') \mapsto \text{Hom}(A \times L, A')$

$\underline{N} = \text{NNO}$ in \underline{A}

WHAT IS STRONG?

STRONG BDN $t: \text{Hom}_{\underline{A}} \Rightarrow \text{Hom}_{\underline{A}}^L$

$$\begin{array}{ccccc}
 A \xrightarrow{\alpha} A & \xrightarrow{t} & A \times L \xrightarrow{t(\alpha)} & A & \\
 X \times A \xrightarrow{X \times \alpha} X \times A & \xrightarrow{\quad} & X \times A \times L \xrightarrow{t(X \times \alpha)} & X \times A & \\
 & & \parallel & \parallel & \\
 & & X \times A \times L \xrightarrow{X \times t(\alpha)} & X \times A &
 \end{array}$$

$$t(X \times \alpha) = X \times t(\alpha)$$

Something special about $\text{Hom}_{\underline{A}}, \text{Hom}_{\underline{A}}^L$ allows this.

CLUE: $E(\underline{A}) = \underline{A}^{\cdot 2} = \text{Cat of endos of } \underline{A}$

$$\begin{array}{ccc}
 \text{BDNs} & t: \text{Hom}_{\underline{A}} \longrightarrow \text{Hom}_{\underline{A}} & \\
 \hline
 E(\underline{A}) & \xrightarrow{T} & E(\underline{A}) \\
 U \downarrow & & \downarrow U \\
 & \underline{A} &
 \end{array}$$

\underline{A} cont. closed with equalizers

$\Rightarrow E(\underline{A})$ enriched in \underline{A} & U is an \underline{A} -funct.

$$t \text{ strong} \Leftrightarrow T \text{ } \underline{A}\text{-functor}$$

$E(\underline{A})$ is tensorred, i.e.

$$E(\underline{A})[(A, \alpha), -] : E(\underline{A}) \longrightarrow \underline{A}$$

has a lax $() \otimes (A, \alpha)$.

$$\frac{X \otimes (A, \alpha) \longrightarrow (B, \beta)}{X \longrightarrow E(\underline{A})[(A, \alpha), (B, \beta)]}$$

$$X \otimes (A, \alpha) = (X \times A, X \times \alpha).$$

t strong $\iff T$ preserves \otimes , i.e.

$$T(X \otimes (A, \alpha)) = X \otimes T(A, \alpha)$$

Each functor of an adjoint pair determines the other. For $E(\underline{A})$ $X \otimes (A, \alpha)$ always exists (regardless of whether \underline{A} is cartesian closed)

So we reformulate everything in terms of \otimes rather than Hom .

Let $\underline{A} = (\underline{V}, \otimes, \mathbb{I}, \alpha, \lambda, \rho)$ Monoidal Cat.

A category \underline{B} is \underline{V} -tensored if there is a left \underline{V} -action on it, i.e.

$$\underline{V} \times \underline{B} \xrightarrow{\otimes} \underline{B}$$

such that

$$\lambda: \mathbb{I} \otimes B \xrightarrow{\cong} B$$

$$\alpha: V_1 \otimes (V_2 \otimes B) \xrightarrow{\cong} (V_1 \otimes V_2) \otimes B$$

with obvious coherence conditions w.r.t. α, λ, ρ of \underline{V} .

$$\text{Ex: } \underline{V}, \underline{V}^{\oplus}, E(\underline{V}) = \underline{V}^{\oplus}$$

If $() \otimes B: \underline{V} \rightarrow \underline{B}$ has adj

$\underline{B}[B, -]: \underline{B} \rightarrow \underline{V}$, then \underline{B} becomes

a \underline{V} -category.

If \underline{B} and \underline{C} are \underline{V} -tensored
 A functor $F: \underline{B} \rightarrow \underline{C}$ is
strong if it comes with "strength"
 morphisms

$$s_{V,B}: V \otimes F(B) \longrightarrow F(V \otimes B)$$

- Natural in V & B
- Compatible with λ, α .

If \underline{B} and \underline{C} are also \underline{V} -cats

$$\begin{array}{c} F: \underline{B} \longrightarrow \underline{C} \quad \text{Strong} \\ \hline F: \underline{B} \longrightarrow \underline{C} \quad \underline{V}\text{-functor} \end{array}$$

(c.f. RJ Wood - thesis)

A nat. trans. $t: F \rightarrow G$ is strong if

$$\begin{array}{ccc} V \otimes F(B) & \xrightarrow{s} & F(V \otimes B) \\ V \otimes t(B) \downarrow & \cong & \downarrow t(V \otimes B) \\ V \otimes G(B) & \xrightarrow{s} & G(V \otimes B) \end{array}$$

$$\frac{t: F \longrightarrow G \text{ strong}}{t: F \longrightarrow G \text{ } \underline{V}\text{-mat}}$$

Have a 2-cat V-Tens

Have strong adjoints

$$\underline{B} \begin{array}{c} \xrightarrow{F} \\ \xleftarrow{U} \end{array} \underline{C}$$

$$F \dashv_{\underline{V}} U \iff F \dashv U \ \& \ F, U \text{ strong} \\ \& \ V \otimes FB \xrightarrow[\cong]{s} F(V \otimes B)$$

equiv:
$$\frac{V \otimes FB \longrightarrow C}{V \otimes B \longrightarrow UC} \cong \text{Nat in } V, B, C.$$

Exi $U: E(\underline{V}) \longrightarrow \underline{V}$ is strong.

U has strong adj \iff

\underline{V} has a NNO (in \otimes sense)

$$\begin{array}{ccccc}
 V \otimes I & \xrightarrow{V \otimes 0} & V \otimes N & \xrightarrow{V \otimes s} & V \otimes N \\
 p \downarrow \cong & & \exists! \downarrow h & & \downarrow h \\
 V & \xrightarrow{g} & A & \xrightarrow{\alpha} & A
 \end{array}$$

$$\frac{h: V \otimes F(I) \longrightarrow (A, \alpha)}{V \otimes I \longrightarrow A} \cong$$

Thm i

$$\frac{\text{Strong BDN} : \text{Hom}_{\underline{V}} \longrightarrow \text{Hom}_{\underline{V}}^L}{L \longrightarrow N}$$

For $F, G: \underline{B}^{\text{op}} \times \underline{B} \longrightarrow \underline{C}$

$t: F \longrightarrow G$ BDN strong if

$$V \otimes t_B = t_{V \otimes B} \quad ?$$

$$\begin{array}{ccc}
 F(B, B) & \xrightarrow{t_B} & G(B, B) \\
 \downarrow V \otimes - & & \downarrow V \otimes - \\
 F(V \otimes B, V \otimes B) & \xrightarrow{t_{V \otimes B}} & G(V \otimes B, V \otimes B)
 \end{array} \quad ?$$

A strong bifunctor $F: \underline{B}^{\text{op}} \times \underline{B} \longrightarrow \underline{C}$

$$\sigma_{V, B, B'} : F(B, B') \longrightarrow F(V \otimes B, V \otimes B')$$

- Nat in B, B' ; Dinat in V
- Compatible with λ, α

$$\begin{array}{ccc}
 F(B, B') & \xrightarrow{\sigma} & F(V \otimes B, V \otimes B') \\
 \downarrow \sigma & & \searrow \sigma \\
 & & F(V' \otimes (V \otimes B), V' \otimes (V \otimes B')) \\
 & & \downarrow F(*, \alpha) \\
 F((V' \otimes V) \otimes B, (V' \otimes V) \otimes B) & \xrightarrow{F(\alpha, *)} & F(V' \otimes (V \otimes B), (V' \otimes V) \otimes B')
 \end{array}$$

B tensored, C no conditions.

Ex: $\text{Hom}_{\underline{A}} : \underline{A}^{\text{op}} \times \underline{A} \longrightarrow \underline{\text{Set}}$

$$\text{Hom}_{\underline{A}}(A, A') \longrightarrow \text{Hom}_{\underline{A}}(V \otimes A, V \otimes A')$$

$$(A \xrightarrow{f} A') \longmapsto (V \otimes A \xrightarrow{V \otimes f} V \otimes A')$$

But not $\underline{A}[-, -] : \underline{A}^{\text{op}} \times \underline{A} \longrightarrow \underline{V}$ if \exists .

$$\underline{A}[A, A'] \xrightarrow{?} \underline{A}[V \otimes A, V \otimes A']$$

$$(W \rightarrow \underline{A}[A, A']) \xrightarrow{?} (W \rightarrow \underline{A}[V \otimes A, V \otimes A'])$$

$$(W \otimes A \rightarrow A') \xrightarrow{??} (W \otimes V \otimes A \rightarrow V \otimes A')$$

$$\text{Hom}^L : \underline{V}^{\text{op}} \times \underline{V} \longrightarrow \underline{\text{Set}}$$

$$\text{Hom}^L(A, B) = \{A \otimes L \rightarrow B\} = \text{Hom}(A \otimes L, B).$$

Examples : ? Profunctors !

SET

function $B \rightarrow C$

relation $B \rightarrow C$

$$B \longrightarrow 2^C$$

$$B \times C \longrightarrow 2$$

$$R \subseteq B \times C$$

(graph)

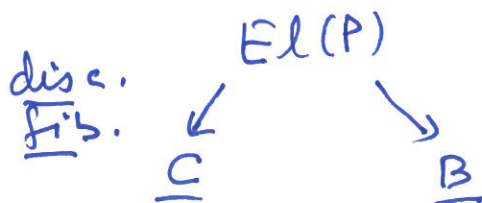
CAT

functor $\underline{B} \rightarrow \underline{C}$

profunctor $\underline{B} \dashrightarrow \underline{C}$

$$\bar{P} : \underline{B} \longrightarrow \underline{\text{Set}}^{\underline{C}^{\text{op}}}$$

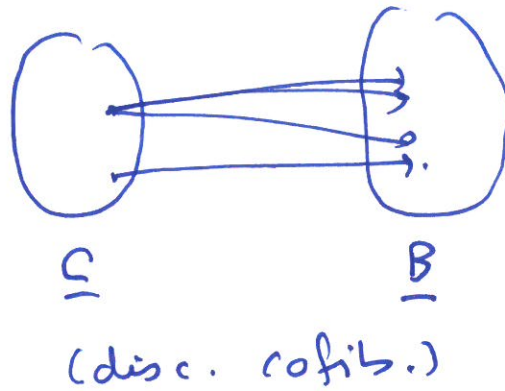
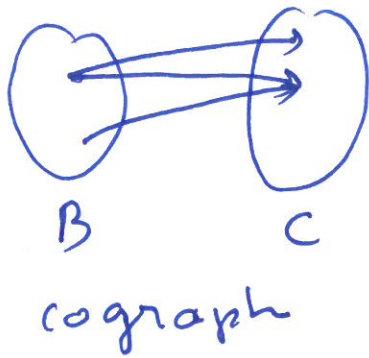
$$P : \underline{C}^{\text{op}} \times \underline{B} \longrightarrow \underline{\text{Set}}$$



$$\begin{bmatrix} 0 & 1 & 1 & 0 & 1 \\ 1 & 1 & 0 & 0 & 0 \\ 1 & 0 & 1 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 \end{bmatrix}$$

Incidence matrix

$$\begin{matrix} \uparrow \\ \leftarrow C \\ \leftarrow \end{matrix} \begin{bmatrix} \cdot & \cdot & \cdot & \cdot \\ \cdot & P(C,B) & \cdot & \cdot \\ \cdot & \cdot & \cdot & \cdot \end{bmatrix} \begin{matrix} \\ \\ \rightarrow B \rightarrow \end{matrix}$$



$$\frac{x \in P(C, B)}{C \xrightarrow[x]{P} B} \text{ denote}$$

$$I_B : \underline{B} \rightarrow \underline{B}$$

$$\text{Hom}_B : \underline{B}^{\text{op}} \times \underline{B} \rightarrow \underline{\text{Set}}$$

$$\underline{B} \xrightarrow{P} \underline{C} \xrightarrow{Q} \underline{D}$$

$$Q \otimes P(D, B) = \int^C Q(D, C) \times P(C, B)$$

$$= \sum_C Q(D, C) \times P(C, B) \sim$$

$$x \otimes y = \left[D \xrightarrow{y} C \xrightarrow{x} B \right]_C$$



Have adjointness for profunctors.

$$F: \underline{B} \rightarrow \underline{C} \text{ functor}$$

$$F_*: \underline{B} \rightarrow \underline{C} \text{ profunctor } F_* = \underline{C}(-, F-)$$

$$F^*: \underline{C} \rightarrow \underline{B}, F^* = \underline{C}(F-, -).$$

The $F_* \dashv F^*$ (The point of it!)



$\underline{B}, \underline{C}$ tensored.

$$F_*: \underline{B} \rightarrow \underline{C} \text{ strong} \iff F: \underline{B} \rightarrow \underline{C} \text{ strong.}$$

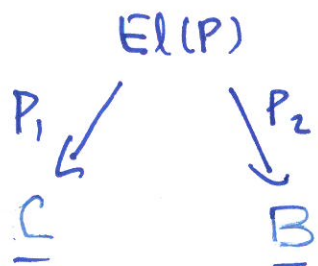
$$F^*: \underline{B} \rightarrow \underline{C} \text{ strong} \iff \exists \alpha: F(V \otimes B) \rightarrow V \otimes F(B) \quad ??$$

$$F_* \dashv F^* \iff V \otimes F(B) \xrightarrow{\cong} F(V \otimes B) !$$

Strong profunctors compose.

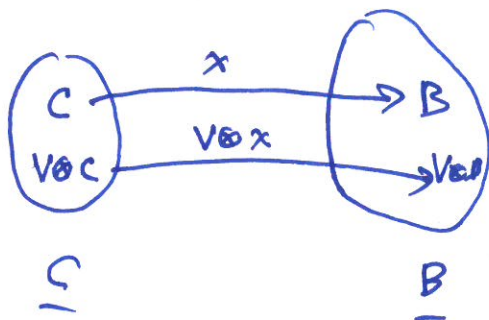
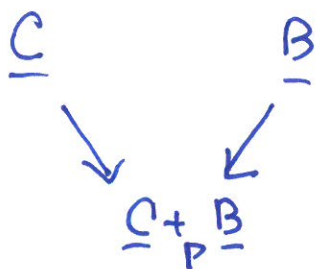
$$\underline{B} \xrightarrow{P} \underline{C}$$

$$\sigma: P(C, B) \longrightarrow P(V \otimes C, V \otimes B)$$



$$x \in P(C, B) \longmapsto V \otimes x \in P(C, B)$$

El(P) is tensored
 P_1, P_2 strong



$\underline{C} + \underline{B}$ tensored

Γ_1, Γ_2 strong

\underline{V} -profunctors $P: \underline{C} \otimes \underline{B} \longrightarrow \underline{V}$

\otimes requires chain in \underline{V} , pres by \otimes .

Relationship ?

For strong profunctors we have a notion of strong natural & dinatural transformations:

$$\begin{array}{ccc}
 P(C, B) & \xrightarrow{t(C, B)} & Q(C, B) \\
 \sigma \downarrow & & \downarrow \sigma \\
 P(V \otimes C, V \otimes B) & \xrightarrow{t(V \otimes C, V \otimes B)} & Q(V \otimes C, V \otimes B)
 \end{array}$$

or $x \in P(C, B)$: $\boxed{V \otimes t(x) = t(V \otimes x)}$