

# Infinite mateship

(title suggested by Mark Weber)

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## Abstract

In this paper, we prove a Cayley Theorem for compact closed categories, in which infinite strings of adjoint endofunctors replace permutations.

## 1 Introduction

A compact closed category may be viewed as a generalised group. Therefore, it makes sense to ask whether there exist a representation theory for compact closed categories.

Unfortunately, I've only got the bare bones of such a theory, without any good combinatorial examples to show it's worthwhile.

## 2 Background

### Definition 2.1

A *compact closed category* is a monoidal category  $(\mathcal{K}, i, \otimes)$  together with an (adjoint) equivalence

$$\mathcal{K}^{\text{op}} \begin{array}{c} \xrightarrow{(-)^{\sharp}} \\ \xleftarrow{(-)^{\flat}} \end{array} \mathcal{K}$$

and di-natural transformations  $\vec{\varepsilon}, \vec{\eta}$  whose components have the form

$$i \xrightarrow{\vec{\eta}_x} x^{\sharp} \otimes x \qquad x \otimes x^{\sharp} \xrightarrow{\vec{\varepsilon}_x} i$$

and which satisfy the so-called *triangle identities*—*i.e.*, that the composites

$$\begin{array}{c} x \xrightarrow{\vec{v}_x^{-1}} x \otimes i \xrightarrow{\iota_x \otimes \vec{\eta}_x} x \otimes (x^{\sharp} \otimes x) \xrightarrow{\alpha_{x, x^{\sharp}, x}} (x \otimes x^{\sharp}) \otimes x \xrightarrow{\vec{\varepsilon}_x \otimes \iota_x} i \otimes x \xrightarrow{\vec{v}_x} x \\ x^{\sharp} \xrightarrow{\vec{v}_{x^{\sharp}}^{-1}} i \otimes x^{\sharp} \xrightarrow{\vec{\eta}_x \otimes \iota_{x^{\sharp}}} (x^{\sharp} \otimes x) \otimes x^{\sharp} \xrightarrow{\alpha_{x^{\sharp}, x, x^{\sharp}}} x^{\sharp} \otimes (x \otimes x^{\sharp}) \xrightarrow{\iota_{x^{\sharp}} \otimes \vec{\varepsilon}_x} x^{\sharp} \otimes i \xrightarrow{\vec{v}_{x^{\sharp}}} x^{\sharp} \end{array}$$

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should equal the identities on  $x$  and  $x^\sharp$  respectively.

A compact closed category is called *strict* if  $(\mathcal{K}, i, \otimes)$  is strict and if the equivalence

$$\mathcal{K}^{\text{op}} \begin{array}{c} \xrightarrow{(-)^\sharp} \\ \xleftarrow{(-)^b} \end{array} \mathcal{K}$$

is an isomorphism.

Note that we do not require a symmetry, or even braiding, on  $\otimes$ ; neither do we require  $(-)^b$  to co-incide with  $(-)^\sharp$ , even up to isomorphism.

**Remark 2.2**

There are also have dinatural transformations  $\overleftarrow{\varepsilon}, \overleftarrow{\eta}$  with components of the form

$$i \xrightarrow{\overleftarrow{\eta}_x} x \otimes x^b \qquad x^b \otimes x \xrightarrow{\overleftarrow{\varepsilon}_x} i$$

defined by

$$\begin{array}{c} i \xrightarrow{\overrightarrow{\eta}_{x^b}} x^{b^\sharp} \otimes x \xrightarrow{\sim} x \otimes x^b \\ x^b \otimes x \xrightarrow{\sim} x^b \otimes x^{b^\sharp} \xrightarrow{\overrightarrow{\varepsilon}_{x^b}} i \end{array}$$

which also satisfy the triangle identities.

**Definition 2.3**

Of *mateship*.

**Definition 2.4**

Of *embedding of compact closed categories*.

### 3 Main Result

**Definitions 3.1**

Let  $\mathcal{K}$  be an arbitrary category. We shall write  $\mathbf{End}^{\pm\infty}(\mathcal{K})$  for the category of endofunctors with infinitely many left and right adjoints; *i.e.*, the objects of  $\mathbf{End}^\infty(\mathcal{K})$  are infinite strings of adjoint endofunctors

$$f = (\cdots \dashv f_{-2} \dashv f_{-1} \dashv f_0 \dashv f_1 \dashv f_2 \dashv \cdots)$$

(with specified units and counits,  $\eta^{(n,n+1)}, \varepsilon^{(n,n+1)}$ ) and an arrow  $f \xrightarrow{\lambda} g$  is defined to be a natural transformation  $f_0 \xrightarrow{\lambda_0} g_0$ .

**Theorem 3.2**

Let  $\mathcal{K}$  be an arbitrary category. Then  $\mathbf{End}^{\pm\infty}(\mathcal{K})$  is a strict compact closed monoidal category.

**Proof.** The tensor product  $\bullet$  is defined by:

$$(f \bullet g)_n = \begin{cases} f_n g_n & \text{if } n \text{ is even} \\ g_n f_n & \text{if } n \text{ is odd} \end{cases}$$

(with the obvious unit and counit maps), and the tensor unit  $i$  by

$$i_n = \text{Id}_{\mathcal{K}}$$

for all  $n$ . It should be clear that these operations make  $\mathbf{End}^{\pm\infty}(\mathcal{K})$  a strict monoidal category

To show that  $\mathbf{End}^{\pm\infty}(\mathcal{K})$  is a (strict) compact closed category, we define  $(-)^{\#}$  and  $(-)^{\flat}$  by shifting on objects and mateship on arrows.

For an object

$$f = (\cdots \dashv f_{-1} \dashv f_0 \dashv f_1 \dashv \cdots)$$

of  $\mathbf{End}^{\pm\infty}(\mathcal{K})$ ,  $f^{\#}$  and  $f^{\flat}$  are defined by

$$f_n^{\#} = f_{n+1} \text{ and } f_n^{\flat} = f_{n-1}$$

respectively.

For an arrow  $f \xrightarrow{\lambda} g$   $\lambda^{\#}$  and  $\lambda^{\flat}$  are defined to be the mates of  $f_0 \xrightarrow{\lambda_0} g_0$ ,

$$g_1 \xrightarrow{\lambda_1} f_1 \text{ and } g_{-1} \xrightarrow{\lambda_{-1}} f_{-1}$$

respectively.

Now the arrows  $i \xrightarrow{\vec{\eta}_f} f^{\#} \bullet f$  and  $f \bullet f^{\#} \xrightarrow{\vec{\varepsilon}_f} i$  defined by the chosen unit and counit maps  $\text{Id}_{\mathcal{K}} \xrightarrow{\eta^{(01)}} f_1 f_0$  and  $f_0 f_1 \xrightarrow{\varepsilon^{(01)}} \iota_{\mathcal{K}}$  satisfy the triangle identities tautologically, and so  $\mathbf{End}^{\pm\infty}(\mathcal{K})$  is a compact closed category. Q.E.D.

**Remark 3.3**

Let  $\mathbf{End}^{\infty}(\mathcal{K})$  denote the category of endofunctors with infinitely many right adjoints—*i.e.*, infinite strings of adjoint endofunctors

$$f = (f_0 \dashv f_1 \dashv f_2 \dashv \cdots).$$

Then both  $\bullet$  and the left internal hom-functor

$$\mathbf{End}^{\pm\infty}(\mathcal{K})^{\text{op}} \times \mathbf{End}^{\pm\infty}(\mathcal{K}) \xrightarrow{-\circ} \mathbf{End}^{\pm\infty}(\mathcal{K})$$

(which, as in any compact closed category, is defined by the formula  $f -\circ h := f^{\#} \bullet h$ ), restrict along the forgetful functor  $\mathbf{End}^{\pm\infty}(\mathcal{K}) \longrightarrow \mathbf{End}^{\infty}(\mathcal{K})$  so that  $\mathbf{End}^{\infty}(\mathcal{K})$ ; is a left-closed monoidal category—and which is not, in general, right-closed.

Dually, the category of endofunctors with infinitely many left adjoints—denoted  $\mathbf{End}^{-\infty}(\mathcal{K})$ —is a right-closed monoidal category which is not, in general, left-closed.

This is where the whole thing actually started, and perhaps explains the somewhat awkward notation.

**Theorem 3.4**

Let  $\mathcal{K}$  be a compact closed category. Then there is an embedding of compact closed categories  $\mathcal{K} \xrightarrow{C} \mathbf{End}^{\pm\infty}(\mathcal{K})$  defined by

$$C(x)_n = \begin{cases} x^{\sharp n} \otimes (-) & \text{if } n > 0 \\ x \otimes (-) & \text{if } n = 0 \\ x^{\flat n} \otimes (-) & \text{if } n < 0 \end{cases}$$

**Proof.** In any monoidal closed category we have  $x \otimes (-) \dashv x \multimap (-)$ , but in a compact closed category we also have  $x \multimap (-) \cong x^{\sharp} \otimes (-)$ .

So we do get an infinite string of adjoint endofunctors

$$(\dots \dashv x^{\flat} \otimes (-) \dashv x \otimes (-) \dashv x^{\sharp} \otimes (-) \dashv \dots)$$

—which is what  $C(x)$  is defined to be.

And it's obviously an "embedding"...

Q.E.D.

## 4 The same again, but with additives

**Theorem 4.1**

If a category  $\mathcal{K}$  has biproducts, then so does  $\mathbf{End}^{\pm\infty}(\mathcal{K})$ . Moreover, if  $\mathcal{K}$  is a compact closed category with biproducts, then the Cayley embedding  $\mathcal{K} \xrightarrow{C} \mathbf{End}^{\pm\infty}(\mathcal{K})$  preserves them.

**Proof.** If  $f$  and  $g$  are objects of  $\mathbf{End}^{\pm\infty}(\mathcal{K})$  then, for each  $n$ , we have composable adjunctions

$$\mathcal{K} \begin{array}{c} \xrightarrow{\Delta} \\ \xleftarrow{\oplus} \end{array} \mathcal{K} \times \mathcal{K} \begin{array}{c} \xrightarrow{(f_{n+1}, g_{n+1})} \\ \xleftarrow{(f_n, g_n)} \end{array} \mathcal{K} \times \mathcal{K} \begin{array}{c} \xrightarrow{\oplus} \\ \xleftarrow{\Delta} \end{array} \mathcal{K}$$

and therefore an infinite string of adjunctions

$$f \oplus g = (\dots \dashv f_{-1} \oplus g_{-1} \dashv f_0 \oplus g_0 \dashv f_1 \oplus g_1 \dashv \dots)$$

(where  $f_n \oplus g_n$  denotes the composite  $\Delta; (f_n, g_n); \oplus$ ).

It is then easy to verify the universal property that  $f \oplus g$  is the biproduct of  $f$  and  $g$ .

Q.E.D.

**Remark 4.2**

If  $\mathcal{K}$  has products and coproducts which do not co-incide, then we get nothing. Given Robin Houston's theorem, this should not be surprising.

Extension to abelian categories?

## 5 Future work

To find an interesting compact closed category  $\mathcal{V}$  (with or without biproducts) which admits an embedding into some  $\mathbf{End}^{\pm\infty}(\mathcal{K})$ , where  $\mathcal{K}$  is demonstrably simpler than  $\mathcal{V}$ . For example, where  $\mathcal{K}$  has smaller cardinality, or does not carry a monoidal structure related to that of  $\mathcal{V}$ .

But note that if  $\mathcal{K}$  is finite/posetal/single-iso-class, then  $\mathbf{End}^{\pm\infty}(\mathcal{K})$  is so too.

So, for example  $\mathbf{Vec}_{fd}$  can only be represented on infinite categories—and ones which admit uncountably many endo-natural transformations of the identity functor! I'm not sure we can really do better than  $\mathbf{Vec}_{fd}$  itself.