## MATH/CSCI 2112, DISCRETE STRUCTURES I, FALL 2005

Handout 3: Natural Deduction for Quantifiers

## Rules.



Forall-elimination $(\forall \mathbf{E})$


Exists-Introduction $(\exists \mathbf{I})$


Exists-Elimination ( $\exists \mathbf{E}$ )


In the above rules:

- in $\forall \mathrm{E}$ and $\exists \mathrm{I}, t$ is any term.
- in $\forall \mathrm{I}$ and $\exists \mathrm{E}, u$ is a fresh variable. Here "fresh" means that this variable does not occur anywhere else in the derivation. It may only occur in the subderivation from lines $m-n$. The " $u$ " that is written between the vertical lines on line $m$ is called a guard - it serves as a reminder that $u$ must be fresh in this subderivation. In particular, this means that no formula containing $u$ can be imported (repeated) into lines $m-n$ from outside lines $m-n$. Also, this means that $u$ cannot occur in the formula $\varphi$ in lines $n$ and $n+1$ of $\exists E$.
- in all rules, $A(u)$ means $S_{u}^{x} A$ (substitution of variable $u$ for free variable $x$ ), and $A(t)$ means $S_{t}^{x} A$ (substitution of term $t$ for free variable $x$ ). In all substitutions for a free variable, you must change the name of any bound variables, if necessary, to avoid capture of variables within a quantifier's scope, if that could occur. It often helps to standardize the variables apart before doing a substitution.


## Examples.



| 1 |  | $\forall x P(x, x)$ |  |
| :---: | :---: | :---: | :---: |
| 2 | $u$ | $\forall x P(x, x)$ | R, 1 |
| 3 |  | $P(u, u)$ | $\forall \mathrm{E}, 2$ |
| 4 |  | $\exists z P(u, z)$ | $\exists \mathrm{I}, 3$ |
| 5 |  | $P(y, z)$ | $\forall \mathrm{I}, 2-4$ |
| 6 | $\forall x P($ | $x) \rightarrow \forall y \exists$ | $\rightarrow \mathrm{I}, 1-5$ |

Non-examples.
Non-example 1

| 1 | $\forall x(A$ | $x) \rightarrow B(x))$ |  |
| :---: | :---: | :---: | :---: |
| 2 | $\exists y A(y)$ |  |  |
| 3 | $u$ | $A(u)$ |  |
| 4 |  | $\forall x(A(x) \rightarrow B(x))$ | R, 1 |
| 5 |  | $A(u) \rightarrow B(u)$ | $\forall \mathrm{E}, 4$ |
| 6 |  | $B(u)$ | $\rightarrow \mathrm{E}, 3,5$ |
| 7 |  |  | ヨE, 2, 3-6 |
| 8 | $\exists y$ A $(1)$ | $) \rightarrow B(u)$ | $\rightarrow \mathrm{I}, 2-7$ |

## Non-example 2

| 1 | $\forall x P(x, x)$ |  |
| :---: | :---: | :---: |
| 2 | $P(u, u)$ | $\forall \mathrm{E}, 1$ |
| 3 | $u \mid P(u, u)$ | R, 2 |
| 4 | $\exists z P(u, z)$ | $\exists \mathrm{I}, 3$ |
| 5 | $\forall y \exists z P(y, z)$ | $\forall \mathrm{I}, 2-4$ |
| 6 | $\forall x P(x, x) \rightarrow \forall y \exists z P(y, z)$ | $\rightarrow \mathrm{I}, 1-5$ |

## Non-example 3

| 1 | $\forall x(A(x) \rightarrow \exists y B(x, y))$ |  |
| :--- | :--- | :--- |
| 2 | $A(y)$ |  |
| 3 | $A(y) \rightarrow \exists y B(y, y)$ | $\forall \mathrm{E}, 1$ |
| 4 | $\exists y B(y, y)$ | $\rightarrow \mathrm{E}, 2,3$ |


| 1 | $\forall x(A(x) \rightarrow \exists y B(x, y))$ |  |
| :--- | :--- | :--- |
| 2 | $A(y)$ |  |
| 3 | $\forall x(A(x) \rightarrow \exists z B(x, z))$ | rename bound variables, 1 |
| 4 | $A(y) \rightarrow \exists z B(y, z)$ | $\forall \mathrm{E}, 1$ |
| 5 | $\exists z B(y, z)$ | $\rightarrow \mathrm{E}, 2,4$ |

WRONG, because $u$ is not fresh in lines 3-6 ( $u$ must not occur in lines 6,7,8).

WRONG, because $u$ is not fresh in lines 3-4 ( $u$ cannot be repeated past the guard from line 2 to line 3).

WRONG, because the substitution in line 3 impropertly captured the variable $y$ in the scope of a quantifier.

CORRECT, because now the variable $y$ does not get captured in the substitution in line 4.

## Problems.

Problem 1 Prove the following in natural deduction:
(a) $Q \rightarrow \forall x P(x) \equiv \forall x(Q \rightarrow P(x))$ - assume that $x$ does not occur in $Q$.
(b) $\sim \exists x P(x) \equiv \forall y \sim P(y)$.
(c) $\forall x P(x) \wedge \forall x Q(x) \equiv \forall x(P(x) \wedge Q(x))$.
(d) $\forall x P(x) \vee \forall x Q(x) \vdash \forall x(P(x) \vee Q(x))$.
(e) $\exists x \forall y P(x, y) \vdash \forall y \exists x P(x, y)$.
(f) $\exists x \forall y P(x, y) \vdash \exists z P(z, z)$.
(g) $\exists x P(x) \vee \exists x Q(x) \equiv \exists x(P(x) \vee Q(x))$.
(h) $\exists x(P(x) \wedge Q(x)) \vdash \exists x P(x) \wedge \exists x Q(x)$.
(i) $\exists x P(x, x) \vdash \exists y \exists z P(y, z)$.
(j) $\forall x(A(x) \rightarrow B(x)) \vdash \exists x \sim B(x) \rightarrow \exists x \sim A(x)$.
$(\mathrm{k}) \sim \exists x(A(x) \wedge B(x)) \equiv \forall x(A(x) \rightarrow \sim B(x))$.
(1) $\exists x \forall y P(x, y, x) \vdash \exists x \forall y \exists z P(x, y, z)$.
$(\mathrm{m}) \vdash \forall x(P(x) \rightarrow \exists y P(y))$.
(n) $\vdash \forall x(\forall y P(y) \rightarrow P(x))$.
(o) $\forall x P(x) \vdash \exists x P(x)$.
(p) $\forall x(A(x) \rightarrow B(x)), \forall y(B(y) \rightarrow C(y)) \vdash \forall z(A(z) \rightarrow C(z))$.
(q) $\exists x A(x), \forall x(A(x) \rightarrow B(x)) \vdash \exists x(A(x) \wedge B(x))$.
(r) $\forall x A(x), \exists x(A(x) \rightarrow B(x)) \vdash \exists x(A(x) \wedge B(x))$.
(s) $\sim \exists x(A(x) \vee B(x)) \equiv \forall x \sim A(x) \wedge \forall x \sim B(x)$.
(t) $\exists x P(x) \rightarrow \forall y Q(y) \equiv \forall x \forall y(P(x) \rightarrow Q(y))$.

Problem 2 Prove the following by natural deduction. Note: each of these problems requires the $\sim \sim$-elimination rule.
(u) $Q \rightarrow \exists x P(x) \equiv \exists x(Q \rightarrow P(x))$ - assume that $x$ does not occur in $Q$.
(v) $\sim \forall x P(x) \equiv \exists y \sim P(y)$.
(w) $\exists x(A(x) \wedge B(x)) \equiv \sim \forall x(A(x) \rightarrow \sim B(x))$.
$(\mathrm{x}) \vdash \exists x(\exists y P(y) \rightarrow P(x))$.
$(\mathrm{y}) \sim \forall x(A(x) \wedge B(x)) \equiv \exists x \sim A(x) \vee \exists x \sim B(x)$.
(z) $\forall x P(x) \rightarrow \exists y Q(y) \equiv \exists x \exists y(P(x) \rightarrow Q(y))$.

Problem 3 Prove the equivalences in Problem 1 (a), (b), (c), (g), (k), (s), (t) and Problem 2 (u), (v), (w), (y), (z) by the laws of statement algebra.

Problem 4 In Problem 1 (d), (e), (f), (h), (i), (j), (l), (o), (p), (q), (r), prove that the converse direction does not hold by giving a counterexample.

