# MATH/CSCI 2112, DISCRETE STRUCTURES I, FALL 2005 Handout 3: Natural Deduction for Quantifiers





In the above rules:

- in  $\forall E$  and  $\exists I, t$  is any term.
- in ∀I and ∃E, u is a *fresh variable*. Here "fresh" means that this variable does not occur anywhere else in the derivation. It may only occur in the subderivation from lines m-n. The "u" that is written between the vertical lines on line m is called a *guard* it serves as a reminder that u must be fresh in this subderivation. In particular, this means that no formula containing u can be imported (repeated) into lines m-n from outside lines m-n. Also, this means that u cannot occur in the formula φ in lines n and n + 1 of ∃E.
- in all rules, A(u) means  $S_u^x A$  (substitution of variable u for free variable x), and A(t) means  $S_t^x A$  (substitution of term t for free variable x). In all substitutions for a free variable, you must change the name of any bound variables, if necessary, to avoid capture of variables within a quantifier's scope, if that could occur. It often helps to standardize the variables apart before doing a substitution.

#### **Examples.**

1	$\forall x$	(A(	$x) \to B(x))$	
2		$\exists y$	A(y)	
3		u	A(u)	
4			$\forall x (A(x) \longrightarrow B(x))$	<b>R</b> , 1
5			$A(u) \longrightarrow B(u)$	$\forall E, 4$
6			B(u)	→E, 3, 5
7			$\exists y  B(y)$	∃I, 6
8		$\exists y$	B(y)	∃E, 2, 3–7
9	$\exists y$	$A(\iota$	$(u) \rightarrow \exists y B(y)$	→I, 2–8

1		$\forall x P(x, x)$		
2		u	$\forall x P(x, x)$	<b>R</b> , 1
3			P(u, u)	$\forall E, 2$
4			$\exists z P(u, z)$	∃I, 3
5		$\forall y$	$\exists z  P(y,z)$	∀I, 2–4
6	$\forall x$	P(x	$(x, x) \longrightarrow \forall y \exists z P(y, z)$	$\rightarrow$ I, 1–5

# Non-examples. Non-example 1

1	$\forall x$	$(A(x) \rightarrow B(x))$	
2		$\exists y A(y)$	
3		u  A(u)	
4		$\forall x (A(x) \to B(x))$	<b>R</b> , 1
5		$A(u) \longrightarrow B(u)$	$\forall E, 4$
6		B(u)	$\rightarrow$ E, 3, 5
7		B(u)	∃E, 2, 3–6
8	$\exists y  A(y) \longrightarrow B(u)$		→I, 2–7

#### Non-example 2

1		$\forall x$	P(x,x)	
2		P(	(u, u)	∀ <b>E</b> , 1
3		u	P(u, u)	<b>R</b> , 2
4			$\exists z P(u, z)$	∃I, 3
5		$\forall y$	$\exists z P(y, z)$	∀I, 2–4
6	$\forall x$	P(x	$(x, x) \longrightarrow \forall y \exists z P(y, z)$	$\rightarrow$ I, 1–5

### Non-example 3

1	$\forall x (A(x) \to \exists y  B(x, y))$	
2	A(y)	
3	$A(y) \longrightarrow \exists y  B(y,y)$	$\forall E, 1$
4	$\exists yB(y,y)$	$\rightarrow$ E, 2, 3

WRONG, because u is not fresh in lines 3–6 (u must not occur in lines 6,7,8).

WRONG, because u is not fresh in lines 3–4 (u cannot be repeated past the guard from line 2 to line 3).

WRONG, because the substitution in line 3 impropertly captured the variable y in the scope of a quantifier.

**CORRECT**, because now the variable y does not get captured in the substitution in line 4.

## Problems.

**Problem 1** Prove the following in natural deduction:

(a) 
$$Q \rightarrow \forall x P(x) \equiv \forall x (Q \rightarrow P(x))$$
 — assume that  $x$  does not occur in  $Q$ .  
(b)  $\sim \exists x P(x) \equiv \forall y \sim P(y)$ .  
(c)  $\forall x P(x) \land \forall x Q(x) \equiv \forall x (P(x) \land Q(x))$ .  
(d)  $\forall x P(x) \lor \forall x Q(x) \vdash \forall x (P(x) \lor Q(x))$ .  
(e)  $\exists x \forall y P(x, y) \vdash \forall y \exists x P(x, y)$ .  
(f)  $\exists x \forall y P(x, y) \vdash \forall y \exists x P(x, y)$ .  
(g)  $\exists x P(x) \lor \exists x Q(x) \equiv \exists x (P(x) \lor Q(x))$ .  
(h)  $\exists x (P(x) \land Q(x)) \vdash \exists x P(x) \land \exists x Q(x)$ .  
(i)  $\exists x P(x, x) \vdash \exists y \exists z P(y, z)$ .  
(j)  $\forall x (A(x) \rightarrow B(x)) \vdash \exists x \sim B(x) \rightarrow \exists x \sim A(x)$ .  
(k)  $\sim \exists x (A(x) \land B(x)) \equiv \forall x (A(x) \rightarrow \sim B(x))$ .  
(l)  $\exists x \forall y P(x, y, x) \vdash \exists x \forall y \exists z P(x, y, z)$ .  
(m)  $\vdash \forall x (P(x) \rightarrow \exists y P(y))$ .  
(n)  $\vdash \forall x (\forall y P(y) \rightarrow P(x))$ .  
(o)  $\forall x P(x) \vdash \exists x P(x)$ .  
(p)  $\forall x (A(x) \rightarrow B(x)), \forall y (B(y) \rightarrow C(y)) \vdash \forall z (A(z) \rightarrow C(z))$ .  
(q)  $\exists x A(x), \forall x (A(x) \rightarrow B(x)) \vdash \exists x (A(x) \land B(x))$ .  
(r)  $\forall x A(x), \exists x (A(x) \rightarrow B(x)) \vdash \exists x (A(x) \land B(x))$ .  
(s)  $\sim \exists x (A(x) \lor B(x)) \equiv \forall x \land y (P(x) \rightarrow Q(y))$ .

**Problem 2** Prove the following by natural deduction. Note: each of these problems requires the  $\sim \sim$ -elimination rule.

(u)  $Q \to \exists x P(x) \equiv \exists x (Q \to P(x))$  — assume that x does not occur in Q.

$$\begin{aligned} &(\mathbf{v}) \sim \forall x \, P(x) \equiv \exists y \sim P(y). \\ &(\mathbf{w}) \ \exists x \, (A(x) \land B(x)) \equiv \sim \forall x \, (A(x) \to \sim B(x)). \\ &(\mathbf{x}) \vdash \exists x (\exists y \, P(y) \to P(x)). \\ &(\mathbf{y}) \ \sim \forall x (A(x) \land B(x)) \equiv \exists x \sim A(x) \lor \exists x \sim B(x). \end{aligned}$$

(z) 
$$\forall x P(x) \rightarrow \exists y Q(y) \equiv \exists x \exists y (P(x) \rightarrow Q(y)).$$

**Problem 3** Prove the equivalences in Problem 1 (a), (b), (c), (g), (k), (s), (t) and Problem 2 (u), (v), (w), (y), (z) by the laws of statement algebra.

**Problem 4** In Problem 1 (d), (e), (f), (h), (i), (j), (l), (o), (p), (q), (r), prove that the converse direction does not hold by giving a counterexample.