

# SPAN AND COSPAN REPRESENTATIONS OF WEAK DOUBLE CATEGORIES

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ABSTRACT. We prove that many important weak double categories can be ‘represented’ by spans, using the basic higher limit of the theory: the tabulator. Dually, representations by cospans via cotabulators are also frequent.

## Introduction

Strict double categories were introduced and studied by C. Ehresmann [Eh1, Eh2], the weak notion in our series [GP1 - GP4]. The strict case extends the more usual (if historically subsequent) notion of 2-category, while the weak one extends bicategories, priorly established by Bénabou [Be]. The extension is made clear in Section 3.

This note is about weak double categories and the (horizontal) *tabulator* of a vertical arrow. The latter is the ‘basic’ higher limit of the theory; in fact the main result of [GP1] says that a weak double category has all (horizontal) double limits if and only if it has: double products, double equalisers and tabulators.

We prove here that the existence of tabulators in a weak double category  $\mathbb{A}$  produces, under suitable hypotheses, a lax functor  $S: \mathbb{A} \rightarrow \text{Span}(\mathbf{C})$  with values in the weak double category of spans over the category  $\mathbf{C}$  of horizontal arrows of  $\mathbb{A}$  (Theorem 6). We say that  $\mathbb{A}$  is *span representable* when this functor  $S$  is horizontally faithful.

Many important weak double categories can be represented in this sense, by spans or - dually - by cospans, via cotabulators.

*Outline.* We begin by a brief review of basic notions on weak double categories, from [GP1, GP2], including the weak double categories of spans and cospans, and the (co)tabulator of a vertical arrow.

Sections 6 and 7 give the main definitions and results recalled above, about (co)span representability. Various weak (or strict) double categories are examined in Sections 8 - 12, proving that many of them are both span and cospan representable. Yet the weak double category  $\text{SpanSet}$ , which is trivially span representable, is *not* cospan representable (Section 8), and  $\text{CospSet}$  behaves in a dual way.

Finally, some common patterns in the previous proofs of representability are analysed in Section 13.

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## 1. Definition

A (strict) *double category*  $\mathbb{A}$  consists of the following structure.

(a) A set  $\text{Ob}\mathbb{A}$  of *objects* of  $\mathbb{A}$ .

(b) *Horizontal morphisms*  $f: X \rightarrow X'$  between the previous objects; they form the category  $\text{Hor}_0\mathbb{A}$  of the *objects and horizontal maps* of  $\mathbb{A}$ , with composition written as  $gf$  and identities  $1_X: X \rightarrow X$ .

(c) *Vertical morphisms*  $u: X \twoheadrightarrow Y$  (often denoted by a dot-marked arrow) between the same objects; they form the category  $\text{Ver}_0\mathbb{A}$  of the *objects and vertical maps* of  $\mathbb{A}$ , with composition written as  $v \bullet u$  (or  $u \otimes v$ , in diagrammatic order) and identities written as  $e_X: X \twoheadrightarrow X$  or  $1_X^\bullet$ .

(d) *Double cells*  $a: (u \xrightarrow{f} v)$  with a *boundary* formed of two vertical arrows  $u, v$  and two horizontal arrows  $f, g$

$$\begin{array}{ccc} X & \xrightarrow{f} & X' \\ u \downarrow & & \downarrow v \\ & a & \\ Y & \xrightarrow{g} & Y' \end{array} \quad (1)$$

Writing  $a: (X \xrightarrow{f} v)$  or  $a: (e \xrightarrow{f} v)$  we mean that  $f = 1_X$  and  $u = e_X$ . The cell  $a$  is also written as  $a: u \rightarrow v$  (with respect to its *horizontal* domain and codomain, which are *vertical* arrows) or as  $a: f \twoheadrightarrow g$  (with respect to its *vertical* domain and codomain, which are *horizontal* arrows).

We refer now to the following diagrams of cells, where the first is called a *consistent matrix*  $\begin{pmatrix} a & b \\ c & d \end{pmatrix}$  of cells

$$\begin{array}{ccccc} X & \xrightarrow{f} & X' & \xrightarrow{f'} & X'' \\ u \downarrow & & \downarrow v & & \downarrow w \\ Y & \xrightarrow{g} & Y' & \xrightarrow{g'} & Y'' \\ u' \downarrow & & \downarrow v' & & \downarrow w' \\ Z & \xrightarrow{h} & Z' & \xrightarrow{h'} & Z'' \end{array} \quad \begin{array}{ccc} X & \xrightarrow{1} & X \\ u \downarrow & & \downarrow u \\ Y & \xrightarrow{1} & Y \end{array} \quad \begin{array}{ccc} X & \xrightarrow{f} & X' \\ e \downarrow & & \downarrow e \\ X & \xrightarrow{f} & X' \end{array} \quad (2)$$

(e) Cells have a *horizontal composition*, consistent with the horizontal composition of arrows and written as  $(a | b): (u \xrightarrow{f'f} w)$ , or  $a|b$ ; this composition gives the category  $\text{Hor}_1\mathbb{A}$  of *vertical arrows and cells*  $a: u \rightarrow v$  of  $\mathbb{A}$ , with identities  $1_u: (u \xrightarrow{1} u)$ .

(f) Cells have also a *vertical composition*, consistent with the vertical composition of arrows and written as  $\begin{pmatrix} a \\ c \end{pmatrix}: (u' \bullet u \xrightarrow{f} v' \bullet v)$ , or  $\frac{a}{c}$ , or  $a \otimes c$ ; this composition gives the category  $\text{Ver}_1\mathbb{A}$  of *horizontal arrows and cells*  $a: f \twoheadrightarrow g$  of  $\mathbb{A}$ , with identities  $e_f = 1_f^\bullet: (e \xrightarrow{f} e)$ .

(g) The two compositions satisfy the *interchange laws* (for binary and zeroary compositions), which means that we have, in diagram (2):

$$\left( \frac{a|b}{c|d} \right) = \left( \frac{a}{c} \middle| \frac{b}{d} \right), \quad \left( \frac{1_u}{1_{u'}} \right) = 1_{u' \bullet u}, \quad (e_f | e_{f'}) = e_{f'f}, \quad 1_{e_X} = e_{1_X}. \quad (3)$$

The first condition says that a consistent matrix  $\begin{pmatrix} a & b \\ c & d \end{pmatrix}$  has a precise *pastings*; the last says that an object  $X$  has an *identity cell*  $\square_X = 1_{e_X} = e_{1_X}$ . The expressions  $(a | f')$  and  $(f | b)$  will stand for  $(a | e_{f'})$  and  $(e_f | b)$ , when this makes sense.

$\mathbb{A}$  is said to be *flat* if every double cell  $a: (u \xrightarrow{f} v)$  is determined by its boundary - namely the arrows  $f, g, u, v$ . A standard example is the double category  $\mathbb{R}\text{elSet}$  of sets, mappings and relations, recalled below in Section 8(c).

## 2. Hints at weak double categories

More generally, in a *weak double category*  $\mathbb{A}$  the horizontal composition behaves categorically (and we still have ordinary categories  $\text{Hor}_0\mathbb{A}$  and  $\text{Hor}_1\mathbb{A}$ ), while the composition of vertical arrows is categorical up to *comparison cells*:

- for a vertical arrow  $u: X \twoheadrightarrow Y$  we have a *left unitor* and a *right unitor*

$$\lambda u: e_X \otimes u \rightarrow u, \quad \rho u: u \otimes e_Y \rightarrow u,$$

- for three consecutive vertical arrows  $u: X \twoheadrightarrow Y$ ,  $v: Y \twoheadrightarrow Z$  and  $w: Z \twoheadrightarrow T$  we have an *associator*

$$\kappa(u, v, w): u \otimes (v \otimes w) \rightarrow (u \otimes v) \otimes w.$$

Interchange holds strictly, as above. The comparison cells are *special* (which means that their horizontal arrows are identities) and horizontally invertible. Moreover they are assumed to be *natural* and *coherent*, in a sense made precise in [GP1], Section 7; after stating naturality with respect to double cells, the coherence axioms are similar to those of bicategories.

$\mathbb{A}$  is said to be *unitary* if the unitors are identities, so that the vertical identities behave as strict units - a constraint which in concrete cases can often be easily met. The terminology of the strict case is extended to the present one, as far as possible.

A *lax (double) functor*  $F: \mathbb{X} \rightarrow \mathbb{A}$  between weak double categories amounts to assigning:

(a) two functors  $\text{Hor}_0 F$  and  $\text{Hor}_1 F$ , consistent with domain and codomain

$$\begin{array}{ccc} \text{Hor}_1 \mathbb{X} & \xrightarrow{\text{Hor}_1 F} & \text{Hor}_1 \mathbb{A} \\ \text{Dom} \downarrow & & \downarrow \text{Dom} \\ \text{Hor}_0 \mathbb{X} & \xrightarrow{\text{Hor}_0 F} & \text{Hor}_0 \mathbb{A} \end{array} \quad \begin{array}{ccc} \text{Hor}_1 \mathbb{X} & \xrightarrow{\text{Hor}_1 F} & \text{Hor}_1 \mathbb{A} \\ \text{Cod} \downarrow & & \downarrow \text{Cod} \\ \text{Hor}_0 \mathbb{X} & \xrightarrow{\text{Hor}_0 F} & \text{Hor}_0 \mathbb{A} \end{array} \quad (4)$$

(b) for any object  $X$  in  $\mathbb{X}$ , a special cell, the *identity comparison* of  $F$

$$\underline{F}(X): e_{FX} \rightarrow Fe_X: FX \dashrightarrow FX,$$

(c) for any vertical composite  $u \otimes v: X \dashrightarrow Y \dashrightarrow Z$  in  $\mathbb{X}$ , a special cell, the *composition comparison* of  $F$

$$\underline{F}(u, v): Fu \otimes Fv \rightarrow F(u \otimes v): FX \dashrightarrow FZ.$$

Again, these comparisons must satisfy axioms of naturality and coherence with the comparisons of  $\mathbb{X}$  and  $\mathbb{A}$  [GP2].

### 3. Dualities

A weak double category has a *horizontal opposite*  $\mathbb{A}^h$  (reversing the horizontal direction) and a *vertical opposite*  $\mathbb{A}^v$  (reversing the vertical direction); a strict structure also has a *transpose*  $\mathbb{A}^t$  (interchanging the horizontal and vertical issues).

The prefix ‘co’, as in *colimit*, *coequaliser* or *colax double functor*, refers to horizontal duality, the main one. Let us note that a weak double category whose horizontal arrows are identities is the same as a *bicategory written in vertical*, i.e. with arrows and weak composition in the vertical direction and strict composition in the horizontal one. This is why *the oplax functors of bicategories correspond here to colax double functors*. (Transposing the theory of double categories, as is done in some papers, would avoid this conflict of terminology, but would produce other conflicts at a more basic level: for instance, colimits in **Set** would become ‘op-limits’ in **RelSet** and **SpanSet**.)

### 4. Spans and cospans

For a category  $\mathbf{C}$  with (a fixed choice of) pullbacks there is a weak double category  $\text{Span}(\mathbf{C})$  of *spans* over  $\mathbf{C}$ , which will play here an important role.

Objects, horizontal arrows and their composition come from  $\mathbf{C}$ , so that  $\text{Hor}_0(\text{Span}\mathbf{C}) = \mathbf{C}$ .

A vertical arrow  $u: X \dashrightarrow Y$  is a span  $u = (u', u'')$ , i.e. a diagram  $X \leftarrow U \rightarrow Y$  in  $\mathbf{C}$ , or equivalently a functor  $u: \mathbb{V} \rightarrow \mathbf{C}$  defined on the formal-span category  $\bullet \leftarrow \bullet \rightarrow \bullet$ . A vertical identity is a pair  $e_X = (1_X, 1_X)$ . A cell  $\sigma: (u \overset{f}{\dashrightarrow} v)$  is a natural transformation  $u \rightarrow v$  of such functors and amounts to the left commutative diagram below

$$\begin{array}{ccc} X & \xrightarrow{f} & X' \\ u' \uparrow & & \uparrow v' \\ U & \xrightarrow{m\sigma} & V \\ u'' \downarrow & & \downarrow v'' \\ Y & \xrightarrow{g} & Y' \end{array} \qquad \begin{array}{ccccc} X & \longrightarrow & X' & \longrightarrow & X'' \\ u' \uparrow & & \uparrow v' & & \uparrow w' \\ U & \xrightarrow{m\sigma} & V & \xrightarrow{m\tau} & W \\ u'' \downarrow & & \downarrow v'' & & \downarrow w'' \\ Y & \longrightarrow & Y' & \longrightarrow & Y'' \end{array} \tag{5}$$

We say that the cell  $\sigma$  is *represented* by its middle arrow  $m\sigma: U \rightarrow V$ , which determines it together with the boundary (the present structure is not flat).

The horizontal composition  $\sigma|\tau$  of  $\sigma$  with a second cell  $\tau: v \rightarrow w$  is a composition of natural transformations, as in the right diagram above; it gives the category  $\text{Hor}_1(\text{Span}\mathbf{C}) = \mathbf{Cat}(\vee, \mathbf{C})$ .

The vertical composition  $u \otimes v$  of spans is computed by (chosen) pullbacks in  $\mathbf{C}$

$$\begin{array}{ccccc}
 X & & Y & & Z \\
 & \swarrow & \nearrow & \swarrow & \nearrow \\
 & U & & V & \\
 & \swarrow & \text{---} & \searrow & \\
 & & W & & 
 \end{array}
 \quad W = U \times_Y V. \quad (6)$$

This is extended to double cells, in the obvious way. For the sake of simplicity *we make*  $\text{Span}(\mathbf{C})$  *unitary*, by adopting the ‘unit constraint’ for pullbacks: the chosen pullback of an identity along any morphism is an identity. The associator  $\kappa$  is determined by the universal property of pullbacks.

Dually, for a category  $\mathbf{C}$  with (a fixed choice of) pushouts there is a unitary weak double category  $\text{Cosp}(\mathbf{C})$  of *cospans* over  $\mathbf{C}$ , that is horizontally dual to  $\text{Span}(\mathbf{C}^{\text{op}})$ . We have now

$$\text{Hor}_0(\text{Cosp}\mathbf{C}) = \mathbf{C}, \quad \text{Hor}_1(\text{Cosp}\mathbf{C}) = \mathbf{Cat}(\wedge, \mathbf{C}), \quad (7)$$

where  $\wedge = \vee^{\text{op}}$  is the formal-cospan category  $\bullet \rightarrow \bullet \leftarrow \bullet$ .

A vertical arrow  $u = (u', u''): \wedge \rightarrow \mathbf{C}$  is now a cospan, i.e. a diagram  $X \rightarrow U \leftarrow Y$  in  $\mathbf{C}$ , and a cell  $\sigma: u \rightarrow v$  is a natural transformation of such functors. Their vertical composition is computed with pushouts in  $\mathbf{C}$ ; again, we generally follow the ‘unit constraint’ for pushouts.

## 5. Tabulators

The (horizontal) *tabulator* of a vertical arrow  $u: X \rightarrow Y$  in the weak double category  $\mathbb{A}$  is an object  $T = \top u$  equipped with a double cell  $t_u: e_T \rightarrow u$

$$\begin{array}{ccc}
 T & \xrightarrow{p} & X \\
 e_T \downarrow & t_u & \downarrow u \\
 T & \xrightarrow{q} & Y
 \end{array}
 \quad
 \begin{array}{ccccc}
 H & \xrightarrow{f} & T & \xrightarrow{p} & X \\
 e \downarrow & e_f & e \downarrow & t_u & \downarrow u \\
 H & \xrightarrow{f} & T & \xrightarrow{q} & Y
 \end{array}
 = h, \quad (8)$$

such that the pair  $(T, t_u: e_T \rightarrow u)$  is a universal arrow from the functor  $e: \text{Hor}_0\mathbb{A} \rightarrow \text{Hor}_1\mathbb{A}$  to the object  $u$  of  $\text{Hor}_1\mathbb{A}$ . Explicitly, this means that for every object  $H$  and every cell  $h: e_H \rightarrow u$  there is a unique horizontal map  $f: H \rightarrow T$  such that  $(e_f | t_u) = h$ , as in the right diagram above.

(In [GP2] we also considered a higher dimensional universal property, which was dropped in later papers and is not used here.) We say that  $\mathbb{A}$  *has tabulators* if all of

them exist, or equivalently if the degeneracy functor  $e: \text{Hor}_0\mathbb{A} \rightarrow \text{Hor}_1\mathbb{A}$  has a right adjoint

$$\top: \text{Hor}_1\mathbb{A} \rightarrow \text{Hor}_0\mathbb{A}, \quad e \dashv \top. \quad (9)$$

In this situation *one can try to represent  $\mathbb{A}$  as a weak double category of spans*, as we shall see below.

Dually  $\mathbb{A}$  *has cotabulators* if the degeneracy functor has a left adjoint

$$\perp: \text{Hor}_1\mathbb{A} \rightarrow \text{Hor}_0\mathbb{A}, \quad \perp \dashv e, \quad (10)$$

so that every vertical arrow  $u: X \twoheadrightarrow Y$  has a cotabulator-object  $\perp u$ , equipped with two horizontal morphisms  $i: X \rightarrow \perp u$ ,  $j: Y \rightarrow \perp u$  and a universal cell  $\iota: (u \begin{smallmatrix} i \\ j \end{smallmatrix} e)$ . This *may* allow representing  $\mathbb{A}$  as a weak double category of cospans.

For a category  $\mathbf{C}$  with pullbacks, the tabulator in  $\text{Span}(\mathbf{C})$  of a span  $u = (u', u'') = (X \leftarrow U \rightarrow Y)$  is its central object  $U$ , with projections  $u', u''$  and the obvious cell  $t_u: e_U \rightarrow u$ . The cotabulator is the pushout of the span in  $\mathbf{C}$ , provided it exists. All this cannot be formulated within the bicategory  $\mathbf{Span}(\mathbf{C})$  (vertically embedded in  $\text{Span}(\mathbf{C})$  as specified in Section 3).

## 6. Theorem and Definition (Span representation)

*We suppose that:*

- (a) *the weak double category  $\mathbb{A}$  has tabulators,*
- (b) *the ordinary category  $\mathbf{C} = \text{Hor}_0(\mathbb{A})$  of objects and horizontal arrows has pullbacks.*

*There is then a canonical lax functor, which is trivial in degree zero*

$$S: \mathbb{A} \rightarrow \text{Span}(\mathbf{C}), \quad \text{Hor}_0(S) = \text{id}\mathbf{C}, \quad (11)$$

*and takes a vertical arrow  $u: X \twoheadrightarrow Y$  of  $\mathbb{A}$  to the span  $Su = (p, q): X \twoheadrightarrow Y$  determined by the tabulator  $\top u$  and its projections  $p: \top u \rightarrow X$ ,  $q: \top u \rightarrow Y$ .*

*The lax functor  $S$  will be called the span representation of  $\mathbb{A}$ .*

*Note.* Related results can be found in Niefield [Ni], for weak double categories with vertical companions and adjoints.

**PROOF.** As in Section 5 we write  $t_u: (e \begin{smallmatrix} p \\ q \end{smallmatrix} u)$  the universal cell of the tabulator  $\top u$ .

The action of  $S$  on a cell  $a$  of  $\mathbb{A}$  is described by the following diagram

$$\begin{array}{ccc}
 X & \xrightarrow{f} & X' \\
 u \downarrow & a & \downarrow v \\
 Y & \xrightarrow{g} & Y'
 \end{array}
 \qquad
 \begin{array}{ccccc}
 & & X & \xrightarrow{f} & X' \\
 & p \nearrow & & p' \nearrow & \\
 \top u & - \top a > & \top v & & \\
 & q \searrow & & t_v \searrow & \\
 & & Y & \xrightarrow{g} & Y' \\
 & & & q' \searrow & \\
 & & & & v \downarrow
 \end{array}
 \quad (12)$$

where the cell  $Sa: Su \rightarrow Sv$  (a morphism of spans) is represented by the coherent morphism  $\top a: \top u \rightarrow \top v$ . The latter is determined by the universal property of the universal cell  $t_v$  of the tabulator  $\top v$

$$(\top a | t_v) = (t_u | a) \quad (p' \cdot \top a = fp, q' \cdot \top a = gq), \quad (13)$$

$$\begin{array}{ccccc} \top u & \xrightarrow{\top a} & \top v & \xrightarrow{p'} & X' \\ e \downarrow & e & \downarrow e & t_v & \downarrow v \\ \top u & \xrightarrow{\top a} & \top v & \xrightarrow{q'} & Y' \end{array} = \begin{array}{ccccc} \top u & \xrightarrow{p} & X & \xrightarrow{f} & X' \\ e \downarrow & t_u & \downarrow e & a & \downarrow v \\ \top u & \xrightarrow{q} & Y & \xrightarrow{g} & Y' \end{array}$$

(In the composition  $(\top a | t_v)$  we write  $\top a$  for  $e_{\top a}$ , as already warned at the end of Section 1.)

To define the laxity comparisons, an object  $X$  of  $\mathbb{A}$  gives a special cell  $\underline{S}(X): e_X \rightarrow S(e_X)$  represented by the morphism

$$k_X: X \rightarrow \top e_X, \quad (k_X | t_{e_X}) = \square_X. \quad (14)$$

For a vertical composite  $w = u \otimes v: X \twoheadrightarrow Y \twoheadrightarrow Z$ , the comparison  $\underline{S}(u, v): Su \otimes Sv \rightarrow Sw$  is represented by the morphism  $k_{uv}$  defined below, where  $P = \top u \times_Y \top v$  is a pullback and  $\sigma = \lambda(e_P)^{-1} = \rho(e_P)^{-1}$

$$k_{uv}: P \rightarrow \top w, \quad (k_{uv} | t_w) = \left( \sigma \mid \frac{r | t_u}{s | t_v} \right), \quad (15)$$

$$\begin{array}{ccccc} P & \xrightarrow{k_{uv}} & \top w & \longrightarrow & X \\ e \downarrow & e & \downarrow e & t_w & \downarrow w \\ P & \xrightarrow{k_{uv}} & \top w & \longrightarrow & Z \end{array} = \begin{array}{ccccccc} P & \xlongequal{\quad} & P & \xrightarrow{r} & \top u & \xrightarrow{p} & X \\ & & e \downarrow & e_r & \downarrow e & t_u & \downarrow u \\ & & & & \top u & \xrightarrow{q} & Y \\ & & \sigma & \nearrow & = & \searrow & \\ & & & & \top v & \xrightarrow{p'} & \\ & & e \downarrow & s & \downarrow e & t_v & \downarrow v \\ P & \xlongequal{\quad} & P & \xrightarrow{s} & \top v & \xrightarrow{q'} & Z \end{array}$$

(Note that one cannot apply interchange to  $(r | t_u) \otimes (s | t_v)$ .) Finally we have to verify the coherence conditions of the comparisons of  $S$  (see [GP2], Section 2.1), and we only check axiom (iii) for the right unitor.

For a vertical map  $u: X \twoheadrightarrow Y$  and  $w = u \otimes e_Y$  we have to verify that the following diagram of morphisms of  $\mathbf{C}$  commutes

$$\begin{array}{ccc} \top u \times_Y Y & \xlongequal{\quad} & \top u \\ (1, k_Y) \downarrow & & \uparrow \top(\rho u) \\ \top u \times_Y \top e_Y & \xrightarrow{k_{ue}} & \top w \end{array} \quad (16)$$

where the pullback  $\top u \times_Y Y$  is realised as  $\top u$ , by the unit constraint, and the morphism  $\top(\rho u)$  is defined by:  $(\top(\rho u) | t_u) = (t_w | \rho u)$ .

Equivalently, by applying the (cancellable) universal cell  $t_u$  and the isocell  $\rho = \rho(e_{\top u})$ , we show that

$$(\rho | (1, k_Y) | k_{ue} | \top(\rho u) | t_u) = (\rho | t_u).$$

In fact we have

$$\begin{aligned} & (\rho | (1, k_Y) | k_{ue} | \top(\rho u) | t_u) = (\rho | (1, k_Y) | k_{ue} | t_w | \rho u) \\ & = \left( \rho | (1, k_Y) | \sigma | \frac{r | t_u}{s | t_e} | \rho u \right) = \left( \rho | \rho^{-1} | \frac{(1, k_Y) | r | t_u}{(1, k_Y) | s | t_e} | \rho u \right) \\ & = \left( \frac{\square \top u | t_u}{e_q | e_k | t_e} | \rho u \right) = \left( \frac{\square \top u | t_u}{e_q | \square Y} | \rho u \right) = \left( \frac{t_u}{e_q} | \rho u \right) = (\rho(e_{\top u}) | t_u). \end{aligned} \quad (17)$$

The fourth, fifth and seventh terms of these computations are represented below, with  $P = \top u \times_Y \top e_Y$  and  $k = k_Y$

$$\begin{array}{ccccccc} \top u & \xrightarrow{(1,k)} & P & \xrightarrow{r} & \top u & \xrightarrow{p} & X \equiv X \\ \downarrow e & & \downarrow e & \searrow e_r & \downarrow e & \searrow t_u & \downarrow u \\ \top u & \xrightarrow{(1,k)} & P & \xrightarrow{r} & \top u & \xrightarrow{q} & Y \\ \downarrow e & & \downarrow e & \searrow s & \downarrow e & \searrow p' & \downarrow e_Y \\ \top u & \xrightarrow{(1,k)} & P & \xrightarrow{s} & \top e_Y & \xrightarrow{q'} & Y \equiv Y \end{array}$$

$\rho u$

$$\begin{array}{ccccccc} \top u & \xrightarrow{p} & X & \equiv & X \\ \downarrow e & \searrow \square \top u & \downarrow e & \searrow t_u & \downarrow u \\ \top u & \xrightarrow{q} & Y & \xrightarrow{\rho u} & Y \\ \downarrow e & \searrow q & \downarrow e & \searrow p' & \downarrow e_Y \\ \top u & \xrightarrow{q} & Y & \xrightarrow{k} & \top e_Y & \xrightarrow{q'} & Y \equiv Y \\ \downarrow e & \searrow e_q & \downarrow e & \searrow e_k & \downarrow e & \searrow t_e & \downarrow e_Y \end{array}$$

■

## 7. Span and cospan representability

Let  $\mathbb{A}$  be a weak double category.

(a) We say that  $\mathbb{A}$  is (horizontally) *span representable* if:

- it has tabulators,
- the ordinary category  $\mathbf{C} = \text{Hor}_0(\mathbb{A})$  has pullbacks,
- the span-representation lax functor  $S: \mathbb{A} \rightarrow \text{Span}(\mathbf{C})$  of (11) is *horizontally faithful*.

The last condition means that the ordinary functors  $\text{Hor}_0 S$  and  $\text{Hor}_1 S$  are faithful. This is trivially true for  $\text{Hor}_0(S) = \text{id}\mathbf{C}$ , and also for  $\text{Hor}_1 S$  when  $\mathbb{A}$  is flat.

(b) By horizontal duality, if  $\mathbb{A}$  has cotabulators and  $\mathbf{C} = \text{Hor}_0(\mathbb{A})$  has pushouts we form a colax functor of *cospan representation*

$$C: \mathbb{A} \rightarrow \text{Cosp}(\mathbf{C}), \quad \text{Hor}_0(C) = \text{id}\mathbf{C}, \quad (18)$$

that takes a vertical arrow  $u: X \twoheadrightarrow Y$  of  $\mathbb{A}$  to the cospan  $Cu = (i, j): X \twoheadrightarrow Y$  formed by the cotabulator  $\perp u$  and its ‘injections’  $i: X \rightarrow \perp u, j: Y \rightarrow \perp u$ .

In this situation we say that  $\mathbb{A}$  is *cospan representable* if this colax functor is horizontally faithful.

## 8. Some basic cases

(a) For a category  $\mathbf{C}$  with pullbacks, the weak double category  $\text{Span}(\mathbf{C})$  is span representable, in a strict sense: the functor  $S: \text{Span}(\mathbf{C}) \rightarrow \text{Span}(\mathbf{C})$  is an isomorphism, and even the identity for the natural choice of the tabulator of a span, namely its central object. Dually, for every category  $\mathbf{C}$  with pushouts,  $\text{Cosp}(\mathbf{C})$  is ‘strictly’ cospan representable.

(b) On the other hand it is easy to see that  $\text{Span}\mathbf{Set}$  is *not* cospan representable, while  $\text{Cosp}\mathbf{Set}$  is *not* span representable. For the first fact we consider a morphism of spans  $\sigma: u \rightarrow u$  represented in the left diagram below, where the objects are cardinal sets ( $0 = \emptyset, 1 = \{0\}, 2 = \{0, 1\}$ )

$$\begin{array}{ccc}
 1 & \longrightarrow & 1 \\
 \uparrow & & \uparrow \\
 2 & \xrightarrow{m\sigma} & 2 \\
 \downarrow & & \downarrow \\
 1 & \longrightarrow & 1
 \end{array}
 \qquad
 \begin{array}{ccc}
 0 & \longrightarrow & 0 \\
 \downarrow & & \downarrow \\
 2 & \xrightarrow{m\sigma} & 2 \\
 \uparrow & & \uparrow \\
 0 & \longrightarrow & 0
 \end{array}
 \qquad (19)$$

All the arrows to 1 are determined but the mapping  $m\sigma: 2 \rightarrow 2$  is arbitrary; the cotabulator pushout is  $\perp u = 1$  and  $\perp\sigma$  does not determine  $\sigma$ .

The second counterexample is shown in the right diagram above, where again  $m\sigma: 2 \rightarrow 2$  is arbitrary, the tabulator pullback is  $\top u = 0$  and  $\top\sigma$  does not detect  $\sigma$ .

Similar counterexamples can be given for any category  $\mathbf{C}$  with finite limits (or colimits) and some object with at least two endomorphisms.

(c) The (strict) double category  $\mathbb{A} = \mathbf{RelSet}$  of sets, mappings and relations [GP1] has  $\text{Hor}_0(\mathbb{A}) = \mathbf{Set}$ , relations for vertical arrows and (flat) double cells given by an inequality

in the ordered category of relations

$$\begin{array}{ccc}
 X & \xrightarrow{f} & X' \\
 \downarrow u & \leq & \downarrow v \\
 Y & \xrightarrow{g} & Y'
 \end{array}
 \quad gu \leq vf. \tag{20}$$

Tabulators and cotabulators exist:  $\top u \subset X \times Y$  is the relation itself, as a subset of  $X \times Y$ , while  $\perp u$  is the pushout of the span  $Su = (X \leftarrow \top u \rightarrow Y)$ , or of any span representing the relation.

Since  $\mathbb{R}\mathbf{Set}$  is flat it is automatically span and cospan representable. The same holds replacing  $\mathbf{Set}$  with any regular category with pushouts.

In the examples below we examine other strict or weak double categories, referring to their definition in [GP1, GP2], briefly reviewed here.

## 9. Representing profunctors

The weak double category  $\mathbb{C}\mathbf{at}$  of *categories, functors and profunctors* was introduced in [GP1], Section 3.1. Objects are small categories, a horizontal arrow is a functor and a vertical arrow is a profunctor  $u: X \rightarrow Y$ , defined as a functor  $u: X^{\text{op}} \times Y \rightarrow \mathbf{Set}$ . A cell  $a: (u \xrightarrow{f} v)$  is a natural transformation  $a: u \rightarrow v(f^{\text{op}} \times g): X^{\text{op}} \times Y \rightarrow \mathbf{Set}$ . Compositions and comparisons are known or easily defined.

The cotabulator  $\perp u = X +_u Y$  of a profunctor  $u: X \rightarrow Y$  is the gluing, or collage, of  $X$  and  $Y$  along  $u$ , with new maps given by  $(\perp u)(x, y) = u(x, y)$  and no maps ‘backwards’; the composition of the new maps with the old ones is defined by the action of  $u$ . The inclusions  $i: X \rightarrow \perp u$  and  $j: Y \rightarrow \perp u$  are obvious, as well as the structural cell  $\iota: (u \xrightarrow{i} e)$

$$\iota: u \rightarrow e_{\perp u}(i^{\text{op}} \times j): X^{\text{op}} \times Y \rightarrow \mathbf{Set}, \quad \iota(x, y): u(x, y) = \perp u(x, y). \tag{21}$$

The tabulator  $\top u$  is the *category of elements* of  $u$ , or Grothendieck construction. It has objects  $(x, y, \lambda)$  with  $x \in \text{Ob}X$ ,  $y \in \text{Ob}Y$ ,  $\lambda \in u(x, y)$  and maps  $(f, g)$  of  $X \times Y$  which form a commutative square in the collage  $X +_u Y$

$$\begin{array}{ll}
 (f, g): (x, y, \lambda) \rightarrow (x', y', \lambda') & (f: x \rightarrow x', g: y \rightarrow y'), \\
 g\lambda = \lambda'f & (u(1_x, g)(\lambda) = u(f, 1_y)(\lambda') \in u(x, y')).
 \end{array} \tag{22}$$

The functors  $p, q$  are obvious, and the structural cell  $\tau = t_u: e_{\top u} \rightarrow u$  is the natural transformation

$$\begin{array}{ll}
 \tau: e_{\top u} \rightarrow u(p^{\text{op}} \times q): (\top u)^{\text{op}} \times \top u \rightarrow \mathbf{Set}, \\
 \tau(x, y, \lambda; x', y', \lambda'): \top u(x, y, \lambda; x', y', \lambda') \rightarrow u(x, y'), & (f, g) \mapsto g\lambda = \lambda'f.
 \end{array} \tag{23}$$

$\mathbb{C}\mathbf{at}$  is easily seen to be span and cospan representable. Indeed, for a cell  $a: (u \xrightarrow{f} v)$ , both the functors  $\top a$  and  $\perp a$  determine every component  $a_{xy}: u(x, y) \rightarrow v(fx, gy)$  of the

natural transformation  $a: u \rightarrow v(f^{\text{op}} \times g): X^{\text{op}} \times Y \rightarrow \mathbf{Set}$

$$\begin{aligned} \top a: \top u &\rightarrow \top v, & \top a(x, y, \lambda) &= (fx, gy, a_{xy}(\lambda)), \\ \perp a: \perp u &\rightarrow \perp v, & \perp a(\lambda: x \rightarrow y) &= a_{xy}(\lambda): fx \rightarrow gy \quad (\lambda \in u(x, y)). \end{aligned} \quad (24)$$

## 10. Representing adjoints.

We prove now that the double category  $\mathbf{AdjCat}$  of (small) *categories, functors and adjunctions*, introduced in [GP1], Section 3.5, is also span and cospan representable.

Again  $\text{Hor}_0(\mathbf{AdjCat}) = \mathbf{Cat}$ . A vertical arrow is now an ordinary adjunction, conventionally directed as the *left* adjoint

$$\begin{aligned} u = (u_\bullet, u^\bullet, \eta, \varepsilon): X \dashrightarrow Y, & \quad (u_\bullet: X \rightarrow Y) \dashv (u^\bullet: Y \rightarrow X), \\ \eta: 1_X \rightarrow u^\bullet u_\bullet, & \quad \varepsilon: u_\bullet u^\bullet \rightarrow 1_Y. \end{aligned} \quad (25)$$

A double cell  $a = (a_\bullet, a^\bullet): u \rightarrow v$  is a pair of mate natural transformations, each of them determining the other via the units and counits of the two adjunctions

$$\begin{aligned} a_\bullet: v_\bullet f &\rightarrow gu_\bullet, & a^\bullet: fu^\bullet &\rightarrow v^\bullet g, \\ a^\bullet &= (fu^\bullet \rightarrow v^\bullet v_\bullet fu^\bullet \rightarrow v^\bullet gu_\bullet u^\bullet \rightarrow v^\bullet g), \\ a_\bullet &= (v_\bullet f \rightarrow v_\bullet fu^\bullet u_\bullet \rightarrow v_\bullet v^\bullet gu_\bullet \rightarrow gu_\bullet). \end{aligned} \quad (26)$$

(a) In  $\mathbf{AdjCat}$  the tabulator  $\top u$  of an adjunction  $u = (u_\bullet, u^\bullet): X \dashrightarrow Y$  is the ‘graph’ of the adjunction, namely the following comma category, equipped with the comma-projections  $p, q$  and an obvious cell  $\tau = t_u: (e_p^p u)$

$$\begin{aligned} \top u = (u_\bullet \downarrow Y) &\cong (X \downarrow u^\bullet), & (x, y; c: u_\bullet x \rightarrow y) &\leftrightarrow (x, y; c': x \rightarrow u^\bullet y), \\ p: \top u &\rightarrow X, & q: \top u &\rightarrow Y, \\ \tau_\bullet: u_\bullet p &\rightarrow q: \top u \rightarrow Y, & \tau_\bullet(x, y; c) &= c: u_\bullet x \rightarrow y. \end{aligned} \quad (27)$$

The tabulator of a cell  $a: (u \xrightarrow{f} v)$ , with components  $a_\bullet x: v_\bullet fx \rightarrow gu_\bullet x$ , is the following functor

$$\top a: \top u \rightarrow \top v, \quad \top a(x, y; c: u_\bullet x \rightarrow y) = (fx, gy; g(c).a_\bullet x: v_\bullet fx \rightarrow gy). \quad (28)$$

This proves that  $\mathbf{AdjCat}$  is *span representable*: in fact the component  $a_\bullet x: v_\bullet fx \rightarrow gu_\bullet x$  is determined by  $\top a(x, u_\bullet x; 1: u_\bullet x \rightarrow u_\bullet x) = (fx, gu_\bullet x; a_\bullet x: v_\bullet fx \rightarrow gu_\bullet x)$ .

(b) In  $\mathbf{AdjCat}$  the cotabulator  $C = \perp u = X +_u Y$  is the category consisting of the disjoint union  $X + Y$ , together with new maps  $\hat{c}_x = (x, y; c: u_\bullet x \rightarrow y)^\wedge \in C(x, y)$  from objects of  $X$  to objects of  $Y$  that are ‘represented’ by objects  $(x, y; c: u_\bullet x \rightarrow y)$  of  $\top u = (u_\bullet \downarrow Y)$ ; the composition of the new maps with old maps  $\varphi \in X(x', x)$ ,  $\psi \in Y(y, y')$  is defined in the obvious way

$$\psi.\hat{c}_x.\varphi = (x', y'; \psi.c.u_\bullet(\varphi): u_\bullet x' \rightarrow u_\bullet x \rightarrow y \rightarrow y')^\wedge. \quad (29)$$

The universal cell  $\iota_\bullet: i \rightarrow ju_\bullet: X \rightarrow \perp u$  is given by  $\iota_\bullet x = (1_{u_\bullet x})^\wedge \in C(x, u_\bullet x)$ .

The cotabulator of a cell  $a: (u \xrightarrow{f} v)$ , with components  $a_\bullet x: v_\bullet f x \rightarrow gu_\bullet x$ , works as  $f$  and  $g$  on the old objects and arrows, as  $\top a$  on the new arrows

$$\perp a: \perp u \rightarrow \perp v, \quad \perp a(x, y; h: u_\bullet x \rightarrow y)^\wedge = (fx, gy; g(h).a_\bullet x: v_\bullet f x \rightarrow gy)^\wedge. \quad (30)$$

This determines  $a_\bullet x$  as above.

## 11. Representing $\mathbb{D}bl$

The strict double category  $\mathbb{D}bl$  of *weak double categories, lax functors and colax functors* is a crucial structure, on which the theory of double adjoints is based. We refer the reader to its introduction in [GP2], Section 2, where the non-obvious point of double cells is dealt with.

We prove now that  $\mathbb{D}bl$  is span representable, *horizontally* and *vertically*.

(a) First, every colax functor  $U: \mathbb{A} \rightarrow \mathbb{B}$  has a *horizontal tabulator*  $(T, P, Q, \tau)$ .

The weak double category  $\mathbb{T} = U \downarrow \mathbb{B}$  is a ‘one-sided’ double comma ([GP2], Section 2.5), with strict projections  $P$  and  $Q$ , which can be used as horizontal *or* vertical arrows. Below the cell  $\eta$  is simply represented by the horizontal transformation  $1_Q: Q \rightarrow Q$  and the tabulator cell  $\tau = t_U$  is linked to the comma-cell  $\pi$  by the unit  $\eta$  and counit  $\varepsilon$  of the *companionship* of  $Q$  with ‘itself’ (see [GP2])

$$\begin{array}{ccccc} \mathbb{T} & \xrightarrow{1} & U \downarrow \mathbb{B} & \xrightarrow{P} & \mathbb{A} \\ \downarrow e & & \downarrow Q & & \downarrow U \\ \mathbb{T} & \xrightarrow{Q} & \mathbb{B} & \xrightarrow{1} & \mathbb{B} \end{array} \quad \tau = (\eta | \pi), \quad \pi = \left( \frac{\tau}{\varepsilon} \right). \quad (31)$$

To be more explicit, the tabulator  $\mathbb{T}$  has objects

$$(A, B, b: UA \rightarrow B), \quad (32)$$

with  $A$  in  $\mathbb{A}$  and  $b$  horizontal in  $\mathbb{B}$ . A horizontal arrow of  $\mathbb{T}$

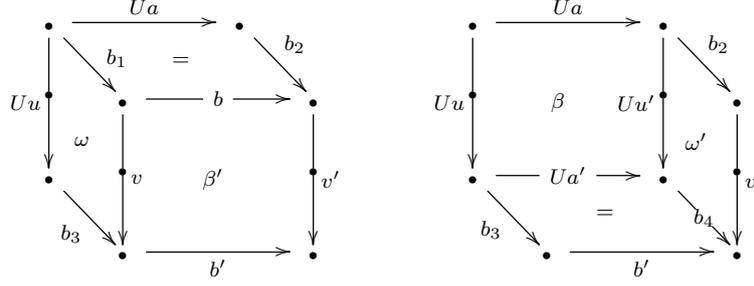
$$(a, b): (A_1, B_1, b_1) \rightarrow (A_2, B_2, b_2), \quad (33)$$

‘is’ a commutative square in  $\text{Hor}_0 \mathbb{B}$ , as in the upper square of diagram (35), below (where the slanting direction must be viewed as horizontal). A vertical arrow of  $\mathbb{T}$

$$(u, v, w): (A_1, B_1, b_1) \rightarrow (A_3, B_3, b_3), \quad (34)$$

‘is’ a double cell in  $\mathbb{B}$ , as in the left square of diagram (35). A double cell  $(\beta, \beta')$  of  $\mathbb{T}$  forms a commutative diagram of double cells of  $\mathbb{B}$

$$(\beta, \beta'): ((u, v, \omega) \xrightarrow{(a, b)} (u', v', \omega')), \quad (\omega | \beta') = (\beta | \omega), \quad (35)$$



The composition laws of  $\mathbb{T}$  are obvious, as well as the (strict) double functors  $P, Q$ . The double cell  $\tau$  has components

$$\tau(A, B, b) = b: UA \rightarrow B, \quad \tau(u, v, \omega) = \omega: Uu \rightarrow v. \quad (36)$$

Its universal property follows trivially from that of the double comma, in [GP2], Theorem 2.6(a).

(b) We have thus a span representation

$$S: \mathbb{Dbl} \rightarrow \text{Span}(\text{LxDbl}), \quad (37)$$

where  $\text{LxDbl} = \text{Hor}_0\mathbb{Dbl}$  is the category of weak double categories and lax functors. (Note that, even though the projections  $P, Q$  of the double comma  $\mathbb{T}$  are strict double functors, a cell  $\varphi: (U \xrightarrow{F} V)$  in  $\mathbb{Dbl}$  gives a lax functor  $\mathbb{T}\varphi: \mathbb{T}U \rightarrow \mathbb{T}V$ .)

To prove that  $\mathbb{Dbl}$  is horizontally span representable, we use the *vertical* universal property of the double comma  $\mathbb{T} = U \downarrow \mathbb{B}$ , in [GP2], Theorem 2.6(b), and deduce the existence of a colax functor  $W: \mathbb{A} \rightarrow \mathbb{T}$  and a cell  $\xi$  such that:

$$\begin{array}{ccc} \mathbb{A} & \xlongequal{\quad} & \mathbb{A} \\ \downarrow W & \xi & \downarrow 1 \\ \mathbb{T}U & \xrightarrow{-P} & \mathbb{A} \\ \downarrow Q & \pi & \downarrow U \\ \mathbb{B} & \xlongequal{\quad} & \mathbb{B} \end{array} = 1_U \quad (QW = U). \quad (38)$$

Now a cell  $\varphi: (U \xrightarrow{F} V)$  in  $\mathbb{Dbl}$  can be recovered from the lax functor  $\mathbb{T}\varphi: \mathbb{T}U \rightarrow \mathbb{T}V$  as follows

$$\begin{aligned} \varphi &= (1_U | \varphi) = (\xi \otimes \pi | e_F \otimes \varphi) = (\xi \otimes t_U \otimes \varepsilon | e_F \otimes \varphi \otimes e_G) \\ &= (\xi | e_F) \otimes (t_U | \varphi) \otimes (\varepsilon | e_G) = (\xi | e_F) \otimes (\mathbb{T}\varphi | t_V) \otimes (\varepsilon | e_G). \end{aligned}$$

(c) Transpose duality leaves  $\mathbb{Dbl}$  invariant up to isomorphism: sending an object  $\mathbb{A}$  to the horizontal opposite  $\mathbb{A}^h$  and transposing double cells we have an isomorphism  $\mathbb{Dbl} \rightarrow \mathbb{Dbl}^t$ . Therefore  $\mathbb{Dbl}$  is also *vertically span representable*, which means that  $\mathbb{Dbl}^t$  is span representable by a lax functor

$$S': \mathbb{Dbl}^t \rightarrow \text{Span}(\text{CxDbI}) \quad (\text{CxDbI} = \text{Hor}_0\mathbb{Dbl}^t = \text{Ver}_0\mathbb{Dbl}). \quad (39)$$

The latter sends a lax functor  $F: \mathbb{A} \rightarrow \mathbb{B}$  to the span  $S'(F) = (\mathbb{A} \leftarrow \mathbb{T}' \rightarrow \mathbb{B})$  associated to its *vertical tabulator*  $(\mathbb{T}, P, Q, \tau)$ , where the weak double category  $\mathbb{T} = \mathbb{B} \downarrow F$  has objects  $(A, B, b: B \rightarrow FA)$ , and the cell  $\tau$  is vertically universal

$$\begin{array}{ccc} \mathbb{T} & \xrightarrow{1} & \mathbb{T} \\ P \downarrow & \tau & \downarrow Q \\ \mathbb{A} & \xrightarrow{F} & \mathbb{B} \end{array} \quad (40)$$

## 12. Theorem (Representing quintets)

*The 2-category  $\mathbf{C}$  is 2-complete if and only if the associated double category  $\mathbb{Q}\mathbf{C}$  of quintets has all double limits. In this case the double category  $\mathbb{Q}\mathbf{C}$  is span representable.*

PROOF. Let us recall that the double category  $\mathbb{Q}\mathbf{C}$  of quintets (introduced by C. Ehresmann) has for horizontal and vertical maps the morphisms of  $\mathbf{C}$ , while its double cells are defined by 2-cells of  $\mathbf{C}$

$$\begin{array}{ccc} X & \xrightarrow{f} & X' \\ u \downarrow & \varphi & \downarrow v \\ Y & \xrightarrow{g} & Y' \end{array} \quad \varphi: vf \rightarrow gu: X \rightarrow Y'. \quad (41)$$

It is known that  $\mathbf{C}$  is 2-complete if and only if it has 2-products, 2-equalisers and cotensors by the arrow-category  $\mathbf{2}$  [St]. First, it is easy to see that 2-products (resp. 2-equalisers) in  $\mathbf{C}$  are ‘the same’ as double products (resp. double equalisers) in  $\mathbb{Q}\mathbf{C}$ . Second, the cotensor  $\mathbf{2} * X$  can be obtained as the tabulator of the vertical identity of  $X$ : they are defined by the same universal property.

Conversely, if the  $\mathbf{C}$ -morphism  $u: X \rightarrow Y$  is viewed as vertical in  $\mathbb{Q}\mathbf{C}$ , its tabulator  $(\top u; p, q; \tau)$  can be constructed as the following inserter  $(\top u; i, \tau)$

$$\top u \xrightarrow{i} X \times Y \xrightleftharpoons[p'']{up'} Y, \quad \tau: up'i \rightarrow p''i: \top u \rightarrow Y, \quad (42)$$

letting  $p = p'i: \top u \rightarrow X$ ,  $q = p''i: \top u \rightarrow Y$  and viewing  $\tau$  as a double cell with boundary  $(1 \begin{smallmatrix} p \\ q \end{smallmatrix} u)$ .

If  $\mathbf{C}$  is 2-complete,  $\mathbb{Q}\mathbf{C}$  is span representable because the lax span representation  $S: \mathbb{Q}\mathbf{C} \rightarrow \text{Span}(\mathbf{C})$  operates on a double cell  $a: vf \rightarrow gu$  of  $\mathbb{Q}\mathbf{C}$  producing a morphism of spans  $Sa: Su \rightarrow Sv$  whose central map  $\top a: \top u \rightarrow \top v$  is defined as follows

$$\begin{array}{ll} t_u: up'i \rightarrow p''i: \top u \rightarrow Y, & t_v: vq'j \rightarrow q''j: \top v \rightarrow Y', \\ j.\top a = (f \times g)i, & t_v.\top a = gt_u.ap'i: vfp'i \rightarrow gp''i. \end{array} \quad (43)$$

$$\begin{array}{ccccc}
\top u & \xrightarrow{i} & X \times Y & \xrightarrow{ap'} & Y \\
\top a \downarrow & & f \times g \downarrow & \xrightarrow{p''} & \downarrow g \\
\top v & \xrightarrow{j} & X' \times Y' & \xrightarrow{q'} & Y'
\end{array}$$

Now  $f$  and  $g$  are determined as the vertical faces of the morphism  $Sa$ . To recover the 2-cell  $a: vf \rightarrow gu$  of  $\mathbf{C}$  from the morphism  $\top a$ , one uses the map  $h: X \rightarrow \top u$  determined by the conditions  $ih = (1, u): X \rightarrow X \times Y$  and  $t_u.h = 1_u$ , so that

$$t_v.\top a.h = (gt_u.ap'i)h = gt_uh.ap'ih = ap'(1, u) = a.$$

■

### 13. Splitting tabulators

In order to ‘explain’ how so many double categories are span representable, we observe that the proof for the non-obvious cases above follows a pattern of the following type (as in Section 12), or a vertical version of the same (as in Section 11). However the argument is rather complicated, and - in the examples above - we preferred to give a direct proof, following this guideline.

We are in a weak double category  $\mathbb{A}$  with tabulators, and the category  $\mathbf{C} = \text{Hor}_0(\mathbb{A})$  has pullbacks. In order that  $\mathbb{A}$  be span representable it is sufficient that, for every vertical arrow  $u$ , there exist two cells  $s_u$  and  $\varepsilon$  satisfying the following condition:

$$\begin{array}{ccccccc}
X & \xrightarrow{h} & \top u & \xrightarrow{p} & X & \xlongequal{\quad} & X \\
\downarrow u & & \downarrow e & & \downarrow u & & \downarrow u \\
& & s_u & & t_u & & \rho u \\
& & \downarrow & & \downarrow & & \\
& & \top u & \xrightarrow{q} & Y & & Y \\
& & \downarrow q_* & & \downarrow e & & \\
Y & \xlongequal{\quad} & Y & \xlongequal{\quad} & Y & \xlongequal{\quad} & Y
\end{array} = 1_u. \quad (44)$$

(Typically,  $q_*$  is the vertical companion of  $q$  and  $\varepsilon = \varepsilon_q$  its counit, but this is not needed in the proof. In the strict case  $\rho u$  is trivial.)

In fact one can recover a cell  $a: (u \xrightarrow{f} v)$  from  $\top a$  (and  $Sa$ ), as follows

$$a = (s_u \mid \frac{t_u}{\varepsilon} \mid \rho u \mid a) = (s_u \mid \frac{t_u}{\varepsilon} \mid \frac{a}{e_g} \mid \rho v) = (s_u \mid \frac{t_u \mid a}{\varepsilon \mid e_g} \mid \rho v) = (s_u \mid \frac{\top a \mid t_v}{\varepsilon \mid e_g} \mid \rho v).$$

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